

6.S084 / 18.S096

LINEAR ALGEBRA AND OPTIMIZATION

TODAY: Monday 9/21

- Pset Due

TODAY 9PM

- Midterm

Wed Oct 7

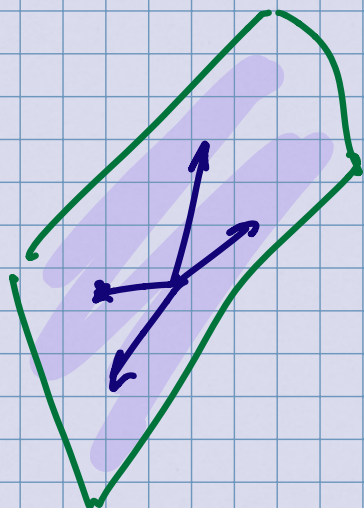
- Check-ups every lecture

- Piazza, Office Hours

Last week

Challenge: Find the dimension of the subspace spanned by

$$\left\{ \begin{pmatrix} 2 \\ 3 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \\ -2 \end{pmatrix}, \begin{pmatrix} 1 \\ 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 0 \\ 4 \end{pmatrix} \right\}$$



4 vectors in $\mathbb{R}^3$

They span a 2-dim subspace!

Can compute, e.g., using Gaussian elim.

Easy proof: All vectors contained
in plane

$$2x - y - z = 0$$

Lecture 8

The RANK of a matrix

Last time:

• Linear independence

$$\left(\begin{array}{c} \sum \lambda_i v_i = 0 \\ \Downarrow \\ \lambda_i = 0 \end{array} \right)$$

+

• Generators

$$\text{span}(v_1, \dots, v_k)$$

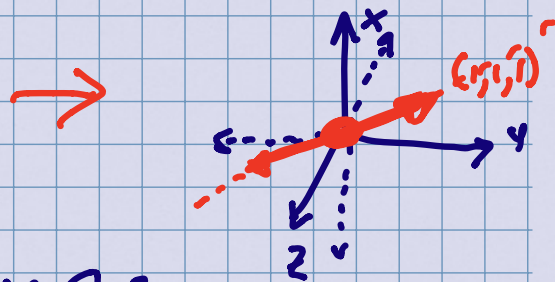
||

• Bases

$$\{v_1, \dots, v_k\}$$

• Dimension

Subspaces:



Two usual descriptions:

• Generators $S = \text{span}\{v_1, \dots, v_k\}$

$\rightarrow S = \text{span}\left\{\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}\right\} \rightarrow S = \text{span}\left\{\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 4 \\ 4 \\ 4 \end{pmatrix}\right\}$

• Equations $S = \{x \mid a_i \cdot x = 0\}$

$\{(x, y, z) \in \mathbb{R}^3 : x = y, y = z, x = z\}$

BASES are very useful.

- But, where do they come from?
- How to compute them?

Claim: If we have generators, can easily "refine" them to a basis.

TASK:

Given subspace $S = \text{span} \{ \underline{v_1, \dots, v_n} \}$
Compute basis B of S .

Algorithm (basic idea)

Initialize basis: $B = \{ \}$

for k from 1 to n $\{v_1, \dots, v_n\}$
if $\underline{v_k}$ is NOT in $\text{span}(B)$
 $B = B \cup \{v_k\}$ (add v_k to basis)
end for
return B .

Claim:

At every step, B is a basis of $\text{span} \{v_1, \dots, v_k\}$.

We already have
an implementation!

Gaussian
elimination

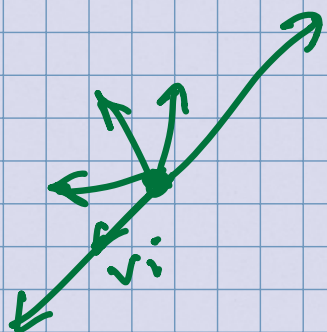


COLUMN RANK : of A .

$$\left\{ \begin{matrix} m \\ A = \end{matrix} \right. \left[\begin{matrix} \overbrace{v_1 \dots v_k}^n \\ \underbrace{\hspace{10em}} \end{matrix} \right]$$

maximum #
of LI columns.

$$0 \leq \text{colrank}(A) \leq n$$



ROW RANK :

$$A = \begin{bmatrix} a_1 \\ \vdots \\ a_m \end{bmatrix}$$

maximum #
of LI rows

$$\begin{cases} a_1 \cdot X = 0 \\ \vdots \\ a_n \cdot X = 0 \end{cases}$$

$$\begin{cases} x=y \\ x=z \\ y=z \end{cases} \rightsquigarrow = \begin{bmatrix} 1 & -1 & 0 \\ -1 & 0 & 1 \\ 0 & 1 & -1 \end{bmatrix}$$

$$0 \leq \text{rowrank}(A) \leq m$$

CLAIM:

FOR ANY MATRIX A

$$\text{ROW RANK}(A) = \text{COLUMN RANK}(A)$$

$$= \boxed{\text{RANK}(A)}$$



Def: The rank of a matrix A
is the # of pivots.

↓ ↓ rank

$$R = EA$$

$$\underline{A} = \begin{bmatrix} \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \end{bmatrix} \xrightarrow{\text{RREF}} R = \begin{bmatrix} \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \end{bmatrix}$$

$$\boxed{\text{rank}(A) = \# \text{ pivots}}$$

and

- pivot columns of A are LI
- pivot columns of R are LI

- pivot rows of R are LI

• A...

(No!)

Example:

$$A = \begin{bmatrix} -1 & 2 & 2 & 1 \\ 0 & 0 & 3 & 2 \\ -2 & 4 & 1 & 0 \end{bmatrix} \xrightarrow{\text{RREF}} R = \begin{bmatrix} 1 & -2 & 0 & 1/3 \\ 0 & 0 & 1 & 2/3 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$E = \begin{pmatrix} 0 & 1/6 & -1/2 \\ 0 & 1/3 & 0 \\ 1 & -1/2 & -1/2 \end{pmatrix}$$

$$\Rightarrow \text{rank}(A) = 2$$

RANK AND FACTORIZATIONS:

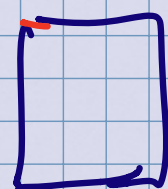
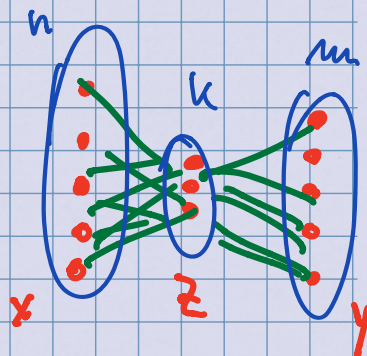
$$\begin{matrix} m \times n \\ \underbrace{\hspace{2cm}} \\ A \end{matrix} = \begin{matrix} m \times k \\ \underbrace{\hspace{2cm}} \\ B \end{matrix} \cdot \begin{matrix} k \times n \\ \underbrace{\hspace{2cm}} \\ C \end{matrix}$$

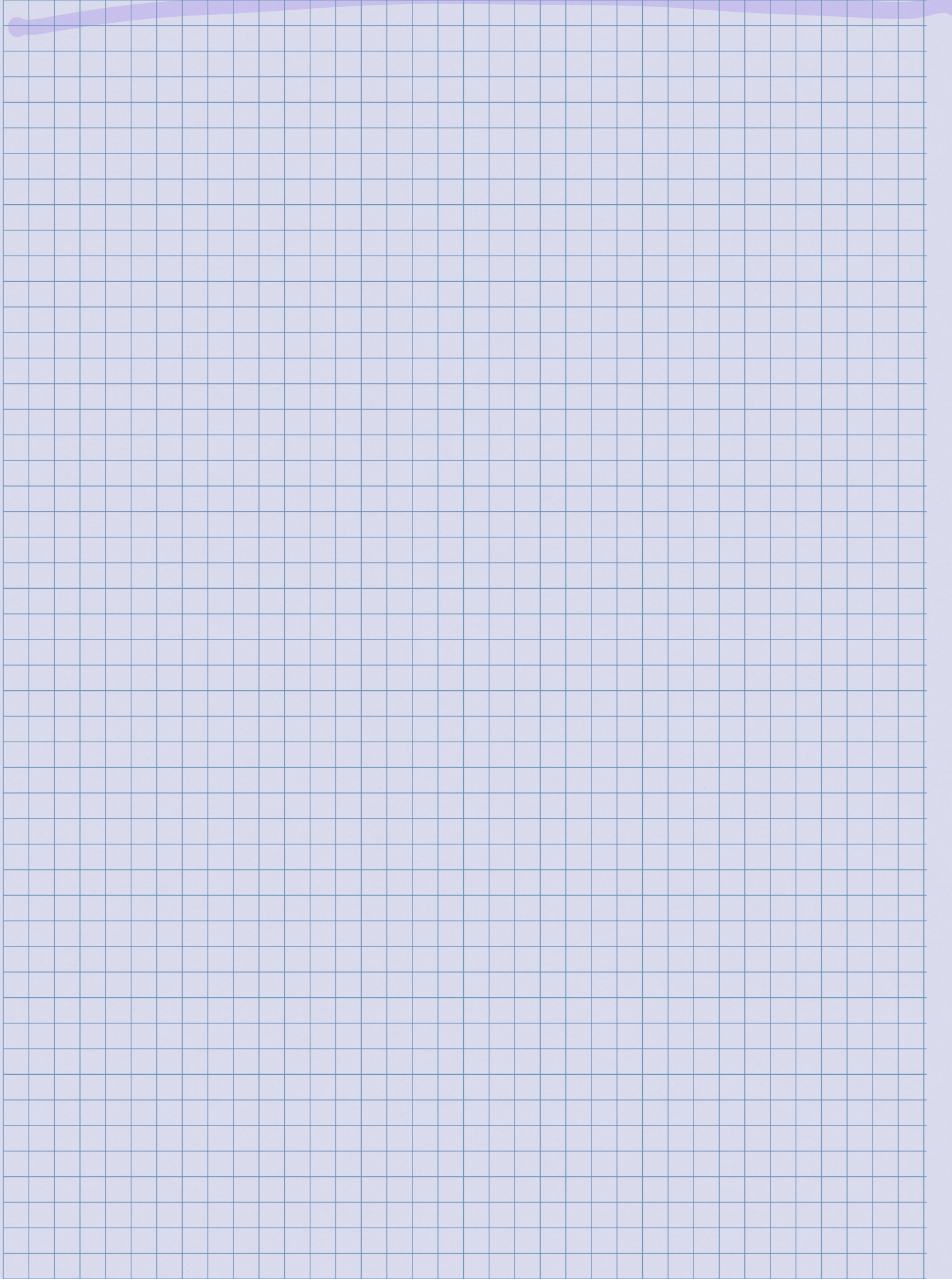
Rank: minimum k for which such factorization exists

Why do we care?

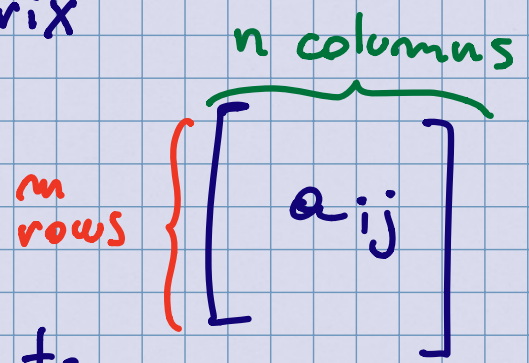
$$y = A \cdot x = \underbrace{B}_{\text{rank } k} \underbrace{(C \cdot x)}_z$$

- Efficient algorithms (e.g. MatMult).
- Dimensionality reduction
- Recommendation systems.
(e.g., Netflix).





① Let A be a $\underline{m} \times \underline{n}$ matrix
($m \geq n$, "tall matrix")



Thm: The following are equivalent:

- The columns of A are linearly independent
- If solvable, $Ax=b$ has a unique solution
- A has a left inverse
- $N(A) = \{0\}$
- $\text{rank}(A) = n$ ("full column rank")

(ii) Let A be a $m \times n$ matrix
($m \leq n$, "fat" matrix)

m rows $\uparrow$

n columns

$\left[\begin{array}{c} a_{ij} \end{array} \right]$

Thm: The following are equivalent:

- The rows of A are linearly independent
- $Ax=b$ is solvable for every b
- A has a right inverse
- $C(A) = \mathbb{R}^m$
- $\text{rank}(A) = m$ ("full row rank")

(iii)

Let A be a square $n \times n$ matrix

$$\left. \begin{array}{c} n \\ \text{rows} \end{array} \right\} \underbrace{\left[a_{ij} \right]}_{n \text{ columns}}$$

Thm: The following are equivalent:

- A is invertible (nonsingular) A^{-1} exists
- $Ax = b$ always has a unique solution
- $N(A) = \{0\}$
- $C(A) = \mathbb{R}^n$
- $\text{rank}(A) = n$

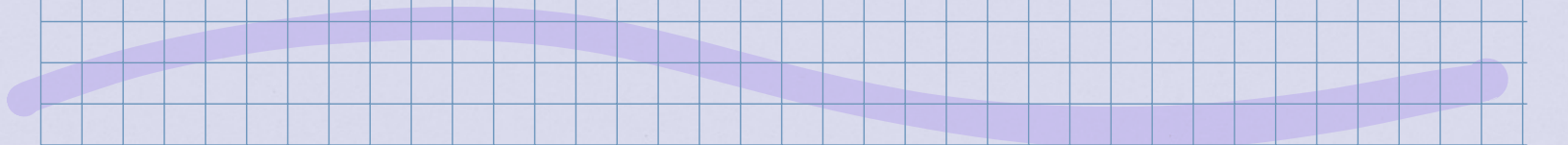
More generally:

$$A \in \mathbb{R}^{m \times n}$$

"Rank-Nullity" Theorem.

$$\text{rank}(A) + \dim N(A) = n$$

(more about this later in the course)





PUZZLE!

Alice "proves":

A has a
left inverse

$\Rightarrow$
(false!)

$Ax = b$ is always
solvable

(NOT A)

"Proof":

$\exists N: NA = I$ (exist
left inv)

$$Ax = b \Rightarrow \overbrace{NA}^I x = Nb \Rightarrow \boxed{x = Nb}$$

Q: What's wrong ???
with this ???